11. Application Servers

CSEP 545 Transaction Processing Philip A. Bernstein

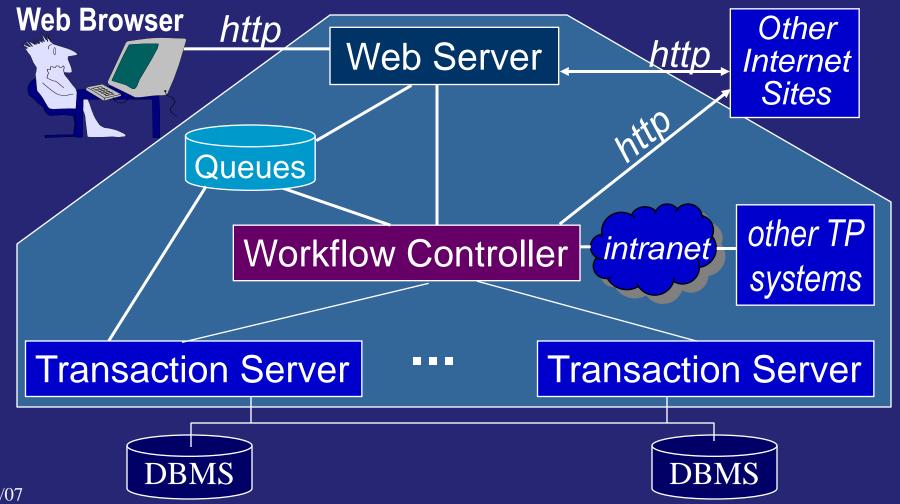
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11.1 Introduction

• An <u>application server</u> coordinates the flow of <u>requests</u> between message sources (displays, applications, etc.) and application programs that run requests as transactions.



Application Server Components

- Web Browser
 - A smart device, with forms, menus, input validation
- Web server
 - Performs front-end work, e.g., security, data caching,
 - "Calls" the web page associated with the URL, which in turn calls a request controller
- Request controller
 - Calls Start, Commit, and Abort
 - App logic that transforms request (automatic loan payment, money transfer) into calls on basic objects (loan, account).
 Sometimes called *business rules*.
- Transaction server
 - Business objects (customer, account, loan, teller)
- DBMS Database Management System

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Application Server Functions

- Glue and veneer for TP applications
 - glue fills in gaps in system functionality
 - covers the interface with a seamless veneer
- Mostly, it provides run-time functions for applications (request control and transaction servers).
 - OS functions: threading and inter-process communication,
 often passed through from the underlying OS.
 - Dist'd system functions: transactions, security, queuing, name service, object pools, load balancing, ...
 - Portal functions: shopping cart, catalog mgmt, personalization ...
- Provides some application development tools
- Provides system mgmt for the running application.

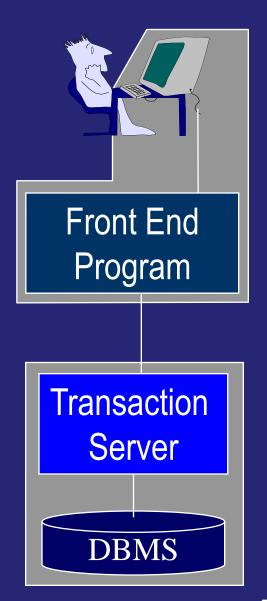
Application Server Products

- Apple WebObjects
- BEA Tuxedo
- BEA Weblogic
- Macromedia ColdFusion
- HP (Tandem) Pathway
- HP (DEC) ACMS
- IBM CICS
- IBM IMS/DC
- IBM Websphere

- Iona iPortal App Server
- Microsoft .NET Enterprise
 Services (formerly COM+,
 MS Transaction Server (MTS))
- Novell SilverStream
- Oracle Application Server
- Sun One Application Server
- Sybase EAServer
- And many others. See serverwatch.internet.com

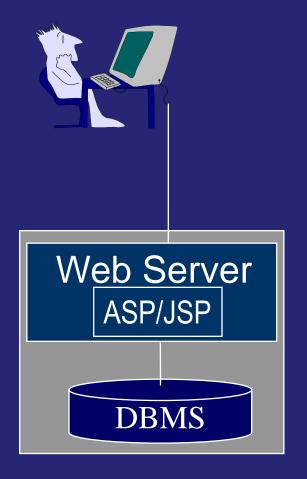
11.2 Two-Tier vs. Three-Tier

- Before the web, most small-to-medium scale apps were implemented in 2 tiers on a LAN
 - PC runs a 4GL, such as Visual Basic,
 PowerBuilder, or Delphi
 - Server system includes transaction server application and DBMS



Two-Tier for the Web

- Front end program \Rightarrow Web server
 - In essence, the web browser is a device
- Web server invokes a web page that has embedded script
 - Active Server Page (ASP .NET) or Java Server Page (JSP)
 - Page (file) extension tells the web server to run the ASP/JSP interpreter
 - Script can include DBMS calls and can run as a transaction



Two-Tier is Enabled by DBMS Stored Procedures

- Stored procedure An application procedure that runs inside the DBMS
 - Often in a proprietary language, such as PL/SQL (Oracle), T-SQL (MS, Sybase)
 - Moving toward standard languages,
 such as Java and C#
- Implement transaction servers as stored procedures
- Use DBMS client-server protocol
- No application server needed
 - Hence, sometimes called "TP lite"

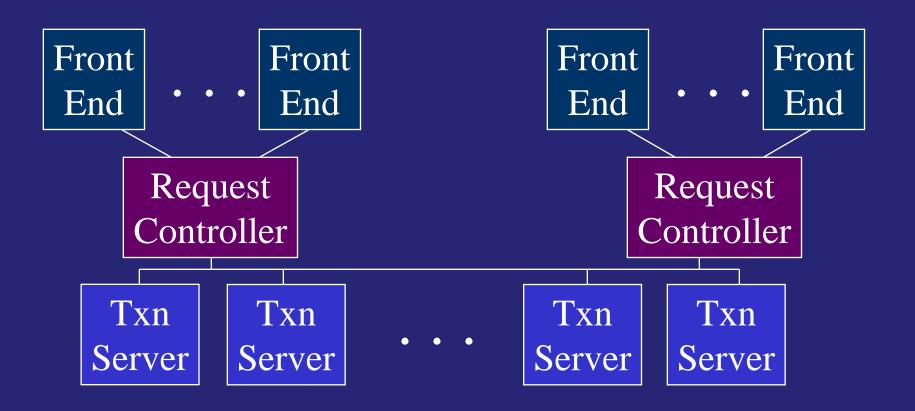
Presentation or Web Server **SQL DBMS** Stored **Procedures SQL** Engine

Scalability Problem of Two-Tier

- 2-tier is feasible, but does not scale as well as 3-tier due to session management
- Session shared state between communicating parties
 - Entails a memory cost and a setup (processing) cost
- Sessions reduce amount of per-request context passing (comm. addresses, authenticated user/device)
 - Standard DB APIs (e.g., ODBC) work this way
 - Hence, in 2-tier, N clients and M servers \Rightarrow N× M sessions
 - E.g. 10^5 presentation servers and 100 servers $\Rightarrow 10^7$ sessions
- Partition presentation servers across workflow controllers

 Each workflow controller still connects to all txn servers but there are many fewer workflow controllers than presentation servers

3-Tier Reduces the Number of Sessions



- Partition the set of front end devices (e.g., 10³ devices per RC)
- $100 \text{ RC} \times (10^3 \text{ devices/RC} + 10^2 \text{ TS/RC}) = 110,000 \text{ sessions}$

Partitioning Txn Servers

- If DB server is a bottleneck, then partition it.
- Request control is needed to direct a call the right DB partition (parameter-based routing)
 - E.g., partition Accounts by account range
 - RC sends a Debit request for Account x to the TS
 connected to the DB partition containing Account x

2-Tier vs. 3 Tier — Other Issues

- In early 90's people argued whether 2-Tier was enough
 - Scalability was the decisive factor, but there were other issues
- Database Servers
 - Nonstandard stored procedure language, usually less expressive with weaker development tools and it's another language to learn
 - Limited interoperability of cross-server calls
 - Limited interoperability of distributed transactions
 - Poor fit with OO design, which are inherently 3-tier (client, business rules, business objects)
- Application Servers
 - more system complexity

How the Web Changed Things

- Front End Program ⇒ Web server
- All requests have to pass through a Web server
 - In 2-tier, each Web server needs sessions to all DB servers
 - Session reduction by request control is less critical but still useful
 - DB partitioning may be implemented by the DB server
- Request control is still useful for request mgmt
 - Calling Start, Commit, and Abort
 - Encapsulating business rules that transform each request into calls on basic objects

11.3 Web Servers

- Presentation independence application is independent of the display device used
 - Today, this is via http and html
 - In the past, it was via a display controller or middle-tier minicomputer whose presentation functions insulated the rest of the back-end system from different device types
- Web server performs presentation functions:
 - Gathering input

– DB caching

Validating input

- Authentication
- They also do some basic request routing

 - Constructing requests
 Invoking applications
- Examples IIS (MS), Apache, Netscape Server

Gathering Input

- Gathering input Select transaction type (menu item, etc.), and fill in a form (request's parameters)
 - Today, Web forms, moving to XML (XForms, XSLT, ...)
- 40 year evolution of presentation devices
 - Teletype, character-at-a-time terminal (async), block-mode terminal (IBM 3270)
 - Specialized devices ATMs, bar code readers, gas pumps, robots, credit card authorization, cash registers, ticket printers, etc.
 - 4GL on a PC ActiveX controls accessed from Visual Basic (VB), PowerBuilder, Delphi, etc.

Caching

- Every process-to-process call has a cost
 - Adds to response time and consumes resources
- Use a cache in Web server to avoid calling workflow controller or DB system
 - Cache popular read-only data that need not be refreshed frequently
 - E.g., catalog items, sale items, cover page at an auction site,
 recent news, etc.
 - Also, data required for input validation info

Input Validation

- Validate input against <u>locally</u> <u>cached</u> tables
 - E.g., product types, department numbers
- Avoids wasting communications and server resources for obvious input errors
 - Fewer round-trips to the DBMS
 - And faster feedback to the end user
- "Cache" is part of the web page
 - List boxes, script
 - Cache size is a factor (it affects page access time)

Authentication

- Authentication determining the identity of a user and/or display device
 - Client system (e.g., PC) may do authentication, but the server usually does it too (doesn't trust clients)
 - Encrypt the wire to avoid wiretapping and spoofing
 (on the Web, use https = http over secure socket layer)
- <u>Geographical entitlement</u> check that a particular *device* is allowed access (e.g., security trading room)
- Need system mgmt functions to create accounts, initialize passwords, bracket hours of access (simplify it using a <u>role</u> abstraction)

Constructing Requests

- A request includes
 - User id for authorization and personalization
 - Device id where to send a reply
 - Device type what message types can it understand?
 - ObjectID in a OO setting
 - RequestID to ask later about request status & to link a reply
 - Request type name of transaction type requested
 - Request-specific parameters

Can be combined with protocol header (e.g., http header)

Application Invocation

- Request arrives as an http message.
 - Need to call a program (i.e. a WFC), to perform the request
- Common Gateway Interface
 - Write a script, store it as a file in cgi-bin
 - Web server creates a process to execute the request (Slow!!)
- ISAPI (Microsoft) and NSAPI (Netscape)
 - Web server calls an in-proc .dll instead of creating a process
 - Web server can cache the .dll
 - More complex programming model, but much faster
- Active Server Pages and Java Server Pages
 - Offers the performance of ISAPI with programmability of CGI

Load Balancing

• Web servers enable <u>scale out</u>, so you can just add more server boxes to handle more load.

- To simplify this problem
 - Ensure all web servers are stateless. I.e., no server-specific state and don't retain client state on web servers (hard to avoid ...)
 - Statelessness implies any web server can process any request.
 - It also makes web server recovery is easy.
 - Randomly assign requests to web servers (e.g., an IP sprayer)

Avoid sending requests to a failed web server

11.4 Transaction Bracketing

- For the most part, Request Controllers (RC) and Transaction Servers are just plain old server programs
- The main RC differentiating features
 - Brackets transactions (issues Start, Commit, and Abort)
 - Handles Aborts (returns cause of the Abort)
 - Does not access the DBMS

Nested Transaction Calls

- What does Start do, when executed within a txn?
 - 1. it starts an independent transaction, or
 - 2. it does nothing, or
 - 3. it increments a nested transaction count (which is decremented by each commit and abort), or
 - 4. it starts a sub-transaction
- (2) and (3) are common.
 - Enables a transaction-bracketed program to be called by another transaction

• (1) implies Be Careful!

Transparent Transaction Bracketing

- Transaction-hood is a property of the app component.
- In COM+, a class is declared:
 - requires new callee always starts a new transaction
 - required if caller is in a transaction, then run callee in caller's transaction, else start a new transaction
 - supported if caller is in a transaction, then run callee in caller's transaction, else run outside of any transaction
 - not supported don't run in a transaction
- Caller can create a transaction context, which supports Commit and Abort (chained model).
 - Callee issues SetComplete when it's done and willing to commit, or SetAbort to abort.

Transparent Txn Bracketing (cont'd)

- Java Enterprise Edition
 - Implements COM+ technology in Java: RequiresNew,
 Required, Supported, NotSupported
 - It came later, so there are two additions.
 - Mandatory If caller is in a transaction, then run the callee in that transaction, else raise an exception
 - Never If caller is in a transaction, then raise an exception

Runtime Library Support

- TP services require runtime library support
 - May or may not be language-specific
- Language-specific
 - Java 2 Enterprise Edition (J2EE, formerly Enterprise Java Beans)
 - Encapsulates runtime library as a container object.
 - BEA Weblogic, IBM Websphere,
 - Older examples are Tandem Pathway (Screen COBOL) and Digital's ACMSxp (Structured Txn Defn Lang)
- Language-independent runtime library
 - MS COM+, IBM's CICS, Oracle App Server, ...

Savepoints

- Savepoint a point in a program where an application saves all its <u>recoverable</u> state
- Can <u>restore</u> a savepoint within the transaction that issued the savepoint. (It's a partial rollback.)
- SQL DBMSs use them to support atomic SQL statements.

```
Start;
get-request;
Savepoint("B"); . . .;
if (error) {Restore("B"); ...; Commit;}
. . .;
Commit;
```

• Savepoints are not recoverable. If the system fails or the transaction aborts, the txn is completely undone.

11.5 Processes and Threads

- Application Server architecture is greatly affected by
 - which components share an address space
 - how many control threads per address space
- TP grew up in the days of batch processing, and reached maturity in the days of timesharing.
- TP users learned early that a process-per-user fails:
 - Too much context switching
 - Too much fixed memory overhead per process
 - Process per user per machine, when distributed
 - Some OS functions scan the list of processes

Load control is hard

Multithreading

- Have multiple threads of control in an address space
- Used to be a major Application Server feature
 - Application Server switches threads when app calls a Application Server function that blocks
- Now, most OS's support it natively
 - Can run a process's threads on different processors (SMP)
- Whether at the user or OS level,
 - multithreading has fewer processes and less context switching
 - but little protection between threads and a server failure affects many transactions

Mapping Servers to Processes

- Presentation/Web servers, workflow controllers, and transaction servers are multithreaded servers
- Costs 1500 25,000 instructions per process call, vs.
 50 instructions per local procedure call ...
 - but it scales, with flexible configuration and control

11.6 Remote Procedure Call

- Program calls remote procedure the same way it would call a local procedure
- Hides certain underlying complexities
 - communications and message ordering errors
 - data representation differences between programs
- Transactional RPC
 - Ideally, Start returns a transaction ID that's hidden from the caller
 - Procedures don't need to explicitly pass transaction id's.

Easier and avoids errors

Binding

Interface definitions

- From app or written in an interface definition language (IDL)
- compiles into Proxy and Stub programs
- Client calls the Proxy (representing the server)
- Stub calls the Server (represents the client on the server)

Marshaling

- proxy <u>marshals</u> (sequentially lays out) calling parameters in a packet and decodes marshaled return values
- stub decodes marshaled calling params and marshals return params

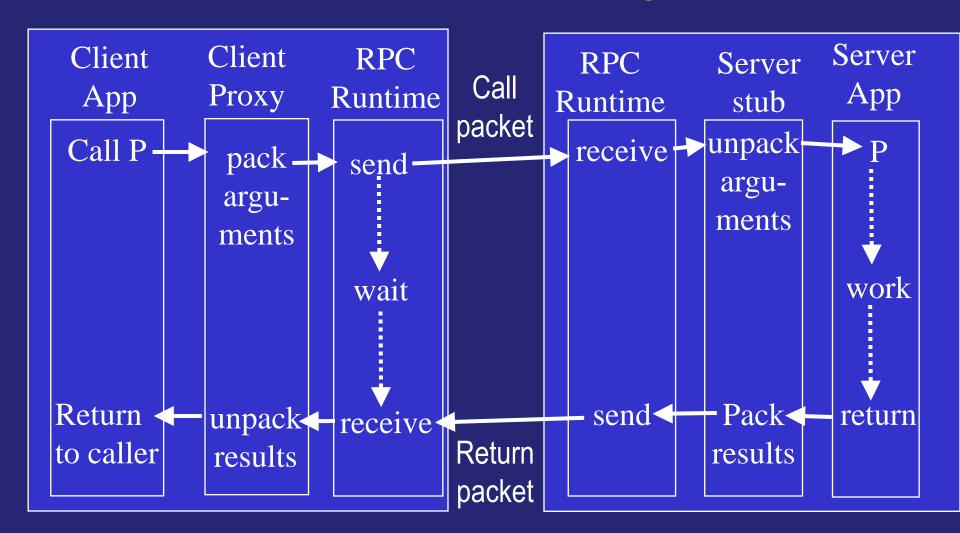
Communications binding

- Client finds the server location via a directory service, based on server name and possibly a parameter value
- To load balance across identical servers, randomly choose a server

Binding (cont'd)

- The binding process has security guarantees
 - The client must have privileges to bind to the server
 - The client must know it's binding to an appropriate server to avoid being spoofed
 - E.g. client and server authenticate each other during session creation, and maybe per-access too

RPC Walkthrough



Client's System

Server's System

Performance

- There are basically 3 costs
 - marshaling and unmarshaling
 - RPC runtime and network protocol
 - physical wire transfer
- In a LAN, these are typically about equal
- Typical commercial numbers are 10-25K machine instructions
- Can do much better in the local case by avoiding a full context switch

Stateful Applications

- Sometimes an application maintains state on client's behalf, possibly across transactions. E.g.,
 - Server scans a file. Each time it hits a relevant record it returns it.
 Next call picks up the scan where it left off.
 - Web server maintains a shopping basket or itinerary, etc.

Approach 1: client passes state to server on each call, and server returns it on each reply. Server retains no state.

- Doesn't work well for TP, because there's too much state
- Note that transaction id context is handled this way.

Stateful Servers (cont'd)

Approach 2: server maintains state, indexed by client id (txn id or cookie). Client's later RPCs must go to same server.

- If the client fails, server must be notified to release client's state or deallocate based on timeout
- For transaction RPC, encapsulate context as a (volatile) resource.
 Delete the state at commit/abort. Or possibly, maintain state across transaction boundaries, but reconstruct it after system failure.
- E.g., COM+: Client can call a server object many times
 - Client creates server object, which retains state across RPCs
 - SetComplete (or SetAbort) by server app says that transaction can be committed (or aborted) and state can be deleted
 - EnableCommit (or <u>DisableCommit</u>) by server app says transaction
 can (or cannot) be committed by client and *don't* delete server state

Fault Tolerance

- If a client doesn't receive a reply within its timeout period
 - RPC runtime can send a "ping" for non-idempotent calls
 - After multiple pings, it return an error.
 - For idempotent calls, RPC runtime can retry the call (server interface definition can say whether it's idempotent)

Web Services

- Distributed computing standards to enable interoperation on the Internet
- SOAP RPC with XML as marshalling format and WSDL as interface definition
- UDDI directory for finding Web Service descriptions
- WS-Transaction 2PC
- WS-Security, WS-Coordination, WS-Routing, ...
- www.ws-i.org

Summary

- Scalability 2 vs. 3 tier, sessions, stored procedures
- Web Server gathering input, validating input, caching, authentication, constructing requests, invoking applications, load balancing
- Transaction bracketing transparency, nesting, exceptions, request integrity, savepoints
- Server processes threads
- RPC binding, stateful servers

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